

- 1) Extensionality:  
 $A = B$  if same elements
- 2) Pairing:  
 $\{a, b\}$
- 3) Union:  
 $A \cup B$
- 4) Infinity:  
 $\omega$
- 5) Power Set:  
 $2^A$
- 6) Replacement:  
 $f: A \rightarrow B$
- 7) Foundation:  
 $\in$
- 8) Choice:  
 $\exists$

Symbols:  
 $\models$  assignment  
 $\wedge \rightarrow \&$   
 $\vee \rightarrow \text{or}$   
 $\neg \rightarrow \text{not}$   
 $\oplus \rightarrow \text{XOR}$

**ORDER OF LOGIC MATTERS:**

$\exists d \in D \forall a \in A. H(a, d)$   
 "There is a single dream all Americans share."

$\forall a \in A \exists d \in D. H(a, d)$   
 "For every American, there is a dream"

validity = true for all values of P(s)

**MAKING CHANGE**

- 1) Prove base cases  
 2) Prove  $P(n+1)$  smallest given  $P(n)$

**NOTES COMPRESSION & CRIB SHEET**

SET NOTATION  
 $e \rightarrow$  element of  
 $\forall \rightarrow$  All  
 $\exists \rightarrow$  There exists  
 $\cup \rightarrow$  union of sets  
 $\cap \rightarrow$  intersection of sets  
 $\bar{\phantom{x}} \rightarrow$  complement  
 $\mathcal{P} \rightarrow$  Power set (set of all subsets)  
 $\{\} \rightarrow$  set  
 $() \rightarrow$  sequence  
 $\lambda \rightarrow$  empty set/sequence  
 $\subset \rightarrow$  subset of

PROPERTIES TO CHECK FOR:  
 • EVENNESS  
 • PRIME NUMBERS

**Binary Relations**

Total  $\rightarrow$  all A have B  
 Surjective  $\rightarrow$  mapped at least once (every B has at least one A)  
 Injective  $\rightarrow$  mapped at most once (at most one A)  
 Bijective  $\rightarrow$  paired

if  $A \text{ surj } B$   
 then  $|A| \geq |B|$

if  $A \text{ inj } B$   
 then  $|A| \leq |B|$

**INDUCTION**

- 1) state base case  
 2) Assume  $P(n)$   
 3) Prove  $P(n+1)$  follows from  $P(n)$

strong induction  
 4) Assume  $P(m)$  men  
 5) Prove  $P(n+1)$  given  $P(i)$  inductive

**PARTIAL ORDERS**

transitivity  
 if  $aRb$  &  $bRc \rightarrow aRc$   
 $\neq$  is transitive  
 $\neq$  is not transitive

symmetry:  
 if  $a \rightarrow b$  then  $b \rightarrow a$   
 strict partial order ( $<$ )  
 is transitive & asymmetric  
 reflexive:  
 if  $a \rightarrow a \forall a$

asymmetry    antisymmetry  
 $\langle \rangle$      $\subseteq$

**RECURSIVE DATA TYPES**

Base Case  
 Constructor Case

$A \text{ exp}$ :  
 substitution Model:  
 $\hookrightarrow$  substitute & evaluate  
 Environment Model:  
 $\hookrightarrow$  evaluate & substitute

**GAMES**

$\hookrightarrow$  use structural induction to prove termination

**MODULAR ARITHMETIC**

main difference is mult. don't cancel (except for w/ primes!)

multiplicative inverse  
 $sp + tk = 1$

**Fermat's Little Theorem:**

$k^p \equiv \text{rem}(k, p)$   
 $k^{p-1} \equiv 1 \pmod{p}$

$\phi(n)$  = # of integers relatively prime to n

$\phi(ab) = \phi(a)\phi(b)$   
 $\phi(p^k) = p^k - p^{k-1}$   
 $k \phi(n) \equiv 1 \pmod{n}$

**PRODUCT ORDER**

$(a, b) (R \times S) (c, d) \leftrightarrow [aRc \wedge bSd]$   
 preserves these, but not totality  
 chain has comparable elements  
 antichains have incomparable elements

**Dilworth's Lemma:**

if largest chain is  $|c|$  then  $c$  antichains can be made  
 $\hookrightarrow$  for all  $t > 0$ , every poset must have a chain of  $t$  or an antichain of at least  $n/t$

Bipartite Graph  
 $\hookrightarrow$  2-colorable

**Hall's theorem:**

Matching of girls and boys iff every subset of girls likes at least as large a set of boys

**NETWORKS**

latency = longest path



**Fund. Theorem**

All  $\mathbb{Z}$  are a product of primes

**SUMS & ASYMPTOTICS**

$H_n = \sum_{i=1}^n \frac{1}{i} \approx \ln(n)$

convert  $\sum \pi \rightarrow \int$  with  $\ln$   
 $n! \approx \left(\frac{n}{e}\right)^n \sqrt{2\pi n}$

**COUNTING**

Pigeonhole Principle

$\binom{n}{k} = \frac{n!}{k!(n-k)!}$

$\binom{n}{k_1, k_2, k_3} = \frac{n!}{k_1! k_2! k_3!}$

BOOKKEEPER RULE

evenness =  $\frac{1}{1-x^2}$   
 $\leq 4 = \frac{1-x^5}{1-x}$

**DIGRAPHS**

DAG,  $D'$  is a strict poset

**STATE MACHINES**

Reach states  $\rightarrow$  check if a preserved invariant is violated (unreachable)  
 partial correctness  $\rightarrow$  final results are always correct  
 termination  $\rightarrow$  the spine eventually stops, always  
 $\hookrightarrow$  look for a desired variable which is strictly decreasing.

**COLORING**

G with max degree  $k$  is  $(k+1)$  colorable

continuous faces (don't touch outside face!)

$v - e + f = 2$  every planar graph has  
 $e \leq 3v - 6$  max(deg)  $\leq 5$   
 $3f \leq 2e$  (5-colorability)

$\frac{1}{m} + \frac{1}{n} = \frac{1}{e} + \frac{1}{2}$  } polyhedra  
 $\uparrow$  #faces at a corner     $\uparrow$  sides per face     $\uparrow$  edges

**PULVERIZER:**

$x \cdot y - r \cdot z = x - y \cdot z$

**Euclid's GCD:**

$(a, b) \rightarrow (b, \text{rem}(a, b))$

Stable Marriage  
 $\hookrightarrow$

**Graphs**

adjacent vertices  $\leftrightarrow$   
 incident edges  $\leftrightarrow$

Handshake Lemma:  
 $2 \text{ deg}(v) = 2e$   
 simple path means no repeat visits  
 simple cycle is a non-repeating loop  
 $k$  connected requires  $k$  edges to be deleted before disconnecting  
 Every Graph has at least  $v - e$  connected components.

**NUMBER THEORY**

- $a|b \rightarrow a|bc \forall c$
- $a|b \ \& \ b|c \rightarrow a|c$
- $a|b \ \& \ a|c \rightarrow a|sb+tc \ \forall s, t$
- $c \neq 0 \ a|b \Leftrightarrow a|cb$

**GCD Properties:**

- Every common divisor divides the GCD
- $\text{GCD}(ka, kb) = k \text{GCD}(a, b)$
- $\text{GCD}(a, b) = 1$  and  $\text{GCD}(a, c) = 1 \rightarrow \text{GCD}(a, bc) = 1$
- $a|bc \ \& \ \text{GCD}(a, b) = 1 \rightarrow a|c$
- $\text{GCD}(a, b) = \text{gcd}(b, \text{rem}(a, b))$

**INCLUSION-EXCLUSION:**

$|S_1 \cup S_2 \cup S_3 \cup \dots \cup S_n|$   
 Sum of sets  
 $-$  sum of 2-way X  
 $+$  sum of 3-way X  
 $-$  sum of ...

**GENERATING FUNCTIONS**

$\langle 1, 1, 1, \dots \rangle = \frac{1}{1-x}$  R-Shift by  $x^k$   
 $\langle 1, -1, 1, -1, \dots \rangle = \frac{1}{1+x}$  AB = convolution  $a * b$   
 $\langle 1, a, a^2, a^3, \dots \rangle = \frac{1}{1-ax}$   
 $\langle 1, 0, 1, 0, \dots \rangle = \frac{1}{1-x^2}$   
 $\langle \binom{k}{0}, \binom{k}{1}, \dots, \binom{k}{k} \rangle = (1+x)^k$

**PROBABILITY**

**4 Step Method:**

- Find Sample Space
- Define Events  $\rightarrow$  Create tree
- Outcome Probabilities

**Birthday Principle:**

$e^{-\frac{1}{d}}$  is max prob of no collisions (approx)